

Duality in Abstract Definability Theory

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Abstract definability theory

Consider:

\mathcal{L} — a **language** (its elements are called **formulas**)

\mathcal{M} — a class of **models** (or **structures**)

\models — a **truth** relation: $M \models A$ between $M \in \mathcal{M}$ and $A \in \mathcal{L}$

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How can we “characterize”

- classes of models from \mathcal{M} definable by a single formula from \mathcal{L} ?
- classes of models from \mathcal{M} definable by a set of formulas from \mathcal{L} ?

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For a set of formulas $\Gamma \subseteq \mathcal{L}$ and a class of models $\mathbb{K} \subseteq \mathcal{M}$, we denote:

$$\begin{aligned}\text{Models}(\Gamma) &:= \{M \in \mathcal{M} \mid M \models \Gamma\} \\ \text{Theory}(\mathbb{K}) &:= \{A \in \mathcal{L} \mid \mathbb{K} \models A\}\end{aligned}$$

The 4 “species” of classes

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Equivalently: if $\mathbb{K} = \bigcap_{i \in I} \mathbb{K}_i$ for some classes $\mathbb{K}_i \in \mathbb{L}$.

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$\mathbb{K} \in \cup \cap \mathbb{L}$ if $\mathbb{K} = \bigcup_{i \in I} \bigcap_{j \in J_i} \mathbb{K}_{i,j}$ for some classes $\mathbb{K}_{i,j} \in \mathbb{L}$.

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$\mathbb{K} \in \mathcal{UL}$ if $\mathbb{K} = \bigcup_{i \in I} \mathbb{K}_i$ for some classes $\mathbb{K}_i \in \mathcal{L}$.

$\mathbb{K} \in \mathcal{UNL}$ if $\mathbb{K} = \bigcup_{i \in I} \bigcap_{j \in J_i} \mathbb{K}_{i,j}$ for some classes $\mathbb{K}_{i,j} \in \mathcal{L}$.

For the “elementary” (i.e. first-order) language \mathcal{L} , the terminology is:

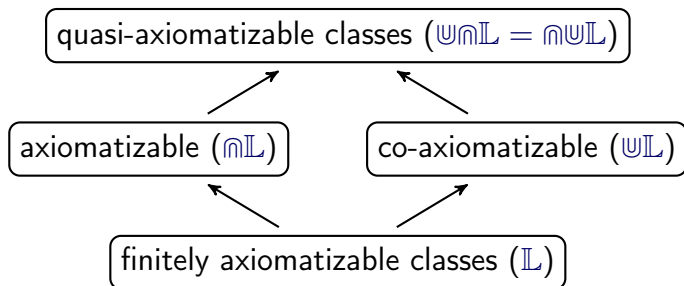
$\mathbb{K} \in \mathcal{L}$ — an **elementary** class of models (or *finitely axiomatizable*)

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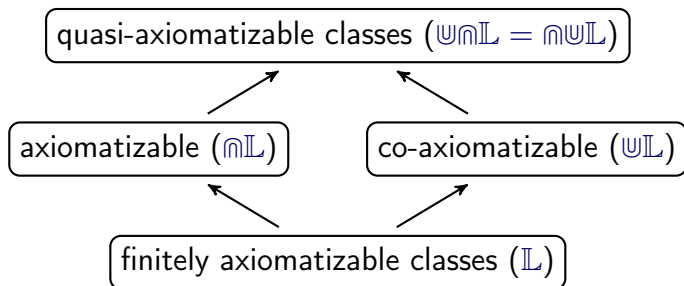
$\mathbb{K} \in \mathcal{UL}$ — a **Σ -elementary** class of models (or *co-axiomatizable?*)

$\mathbb{K} \in \mathcal{UNL}$ — a **$\Sigma\Delta$ -elementary** class of models (or *quasi-axiomatizable?*)

The hierarchy of the 4 species of classes of models



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- Classes in \mathbb{L} : the classes of all groups, all rings, all fields, all orders
- Classes in $\cap\mathbb{L}$: infinite groups, infinite rings, infinite fields, inf. orders
- Classes in $\cup\mathbb{L}$: finite groups, finite rings, finite fields, finite orders
- Classes in $\cap\cup\mathbb{L}$: infinite fields of characteristic $p > 0$;
infinite finitely dimensional vector spaces
- Not even in $\cup\cap\mathbb{L}$: well-ordered sets, periodic groups, simple groups

Characterization of quasi-axiomatizable classes of models

Definition

For models $M, N \in \mathcal{M}$, we write $M \sqsubseteq N$ if,

for every formula $A \in \mathcal{L}$, $(M \models A \implies N \models A)$.

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Theorem (E.Z., 2017)

For any class of models $\mathbb{K} \subseteq \mathcal{M}$, the following conditions are equivalent:

- (A1) \mathbb{K} is the union of intersections of finitely ax. classes: $\mathbb{K} \in \cup \cap \mathcal{L}$
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$$(C1) \quad \mathbb{K} = \bigcup_{M \in \mathbb{K}} \bigcap_{M \models A} \text{Models}(A)$$

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Characterization of quasi-axiomatizable classes (improved)

Notation: all theories of models M from the class \mathbb{K} :

$$\text{TH}(\mathbb{K}) := \{ \text{Theory}(M) \mid M \in \mathbb{K} \} \quad \text{— this is a set!}$$

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Duality of notions in model theory

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Definition. The truth relation \models is said to be of **finite index** if

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Theorem. Conditions (1) and (2) are equivalent!

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- equivalently:
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Examples:

First-order logic, group theory, modal logic \mathbf{K} — are **theories**

True arithmetic — $\text{Theory}(\mathbb{N}, +, \times, =)$ — is a **basic theory**

Characterization of quasi-theories

By analogy with $M \sqsubseteq N$, we introduce, for formulas $A, B \in \mathcal{L}$, the notion:

$A \Rightarrow B$ if for all models M ($M \models A \implies M \models B$).

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As before, we can avoid unions and intersection over *classes*.

Compactness

A class of models $\mathbb{K} \subseteq \mathcal{M}$ is called **compact** if,
for every set of formulas $\Gamma \subseteq \mathcal{L}$, we have:

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Examples of compact classes:

- all groups, linear orders, Boolean algebras, etc.
- all infinite groups, linear orders, Boolean algebras, etc.
- all reflexive Kripke frames or models, etc.

Examples of non-compact class: all finite groups; all well-orders, etc.

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Definition (Truth and Boolean connectives)

- The truth relation \models **respects negation** if

$$M \not\models A \quad \iff \quad M \models \neg A$$

- The truth relation \models **respects conjunction** if

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Remark. If \models respects negation, then $M \subseteq N$ is equivalent to $M \equiv N$.

Characterization of axiomatizable classes

Theorem (E.Z., 2017)

Suppose that the system $(\mathcal{M}, \models, \mathcal{L})$ is such that

- the truth relation \models respects negation and conjunction;
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- \mathbb{K} is *axiomatizable* $\iff \mathbb{K}$ is closed under \equiv and compact
- \mathbb{K} is *co-axiomatizable* $\iff \overline{\mathbb{K}}$ is closed under \equiv and compact
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	Both	\mathbb{K} is ...	$\overline{\mathbb{K}}$ is ...
$\mathbb{K} \in \cup \cap \mathcal{L}$ (\mathbb{K} is quasi-axiomatizable)	\equiv		
$\mathbb{K} \in \cup \mathcal{L}$ (\mathbb{K} is co-axiomatizable)	\equiv		compact
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$\mathbb{K} \in \mathcal{L}$ (\mathbb{K} is finitely axiomatizable)	\equiv	compact	compact

Conclusion and further directions

Conclusions:

- In the relation $M \models A$ there is a natural “symmetry” between the objects on the left (**models**) and the objects on the right (**formulas**).
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More technically:

- What is the notion dual to that of a **compact** class of models?
- What are criteria for $T \subseteq \mathcal{L}$ to be a **theory**, **co-theory**, **basic theory**?

Conclusion and further directions

Conclusions:

- In the relation $M \models A$ there is a natural “symmetry” between the objects on the left (**models**) and the objects on the right (**formulas**).
- One can introduce notions and obtain results that are dual to the known ones: exchange the role of models and formulas (if possible).
- But there are subtleties:
 - usually models form **a class**,
 - but formulas form **a set**.
- Therefore, the main (philosophical) question is: *For which notions and results, the dual notions and results are meaningful and interesting?*

More technically:

- What is the notion dual to that of a **compact** class of models?
- What are criteria for $T \subseteq \mathcal{L}$ to be a **theory**, **co-theory**, **basic theory**?

Thank you!